

PRIMORIAL BASED PRIMES

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(Abstract)

The product of all successive primes starting at Prime[1] = 2 and ending at the mth prime, Prime[m], is called the mth primorial. In order to form primorial based primes an odd number must be added or subtracted from a primorial. This note explores the nature of the primes formed in this process.

The mth primorial is defined as

$$\text{primorial}[m] = \prod_{n=1}^m \text{Prime}[n]$$

where Prime[n] is the nth prime number. The primorial is always an even number so that to produce a prime requires adding or subtracting 1 from a primorial or adding and subtracting an odd number from a primorial. Only the latter cases will be considered here. Some obvious properties of such numbers are :

(1.) $\text{primorial}[m] \pm (2 * k + 1)$ cannot be prime if $(2 * k + 1)$ is less than Prime[m + 1].

(2.) If $(2 * k + 1)$ is greater than Prime[m] but less than Prime[m] ^ 2 then $\text{primorial}[m] \pm (2 * k + 1)$ can be a prime only if $2 * k + 1$ is also a prime.

(3.) For numbers larger than $\text{primorial}[m] \pm \text{Prime}[m]^2$ primes of the form $\text{primorial}[m] \pm (2 * k + 1)$ can be prime even if $(2 * k + 1)$ is not a prime.

It is the purpose of this note to point out the properties and advantages of searching for primes of the numbers

$$N[m, p] = \text{primorial}[m] \pm \text{Prime}[m + p] \quad (1.)$$

where p varies from 1 to p (max) with p (max) the solution to the equation

$$\text{Prime}[m + p (\text{max})] = \text{Prime}[m]^2. \quad (2.)$$

The principle argument for searching for primes of this form is that primorials with more than a million digits, $m = 170\,000$ for example, can be calculated in a matter of a few seconds with most desk top computers. Hence it is of interest to explore some of the properties of such numbers. In what follows searches for primes in the region $N[m, p]$ with p varied from 1 to p (max) ,

labeled region 1, will be compared to the numbers

$$N[m, k] = \text{primorial}[m] \pm (\text{Prime}[m]^2 + 2 * k), \quad (3.)$$

labeled region 2, where k varies from 1 to $\text{Prime}[m] (\text{Prime}[m] - 1) / 2$. Which is the same number of searches required for region 1 if values of $(2 * k + 1)$ that are not prime are also included in the search.

In order to study the properties of the numbers defined in Eqs. 2 and 3 it is convenient to use the "Mathematica" language where $p(\text{max})[m]$ is defined as

$$p(\text{max})[m] = \text{PrimePi}[\text{Prime}[m]^2] - m, \quad (4.)$$

primorial[m] by,

$$\text{primorial}[m] = \text{Product}[\text{Prime}[k], \{k, 1, m\}] \quad (5.)$$

and the number of digits in the primorial[m] which are given as

$$\text{numberdigits}[m] = \text{Sum}[\text{DigitCount}[\text{primorial}[m], 10, j], \{j, 0, 9\}]. \quad (6.)$$

It should be noted that for large primes any number of the form given in Eq. 1 will almost always have the same number of digits as those in the primorial since a prime with 7 to 10 digits will be added or subtracted from a primorial with more than a million digits. Nearly all of the most significant nearly million digits would have to be 9's in the case of adding a prime to a primorial or 1 followed by almost a million zeros in the case of subtracting a prime in order to change the number of digits from that of the primorial.

In order to test for probable primality the following four methods are available in "Mathematica"

$$\begin{aligned} \text{PrimeQ}[N[m, p]] &= \\ \text{True} & \quad (7.) \end{aligned}$$

$$\text{NextPrime}[N[m, p]] = \text{value of } m \text{ for the next larger prime} \quad (8.)$$

$$\text{CompositeQ}[N[m, p]] = \text{False} \quad (9.)$$

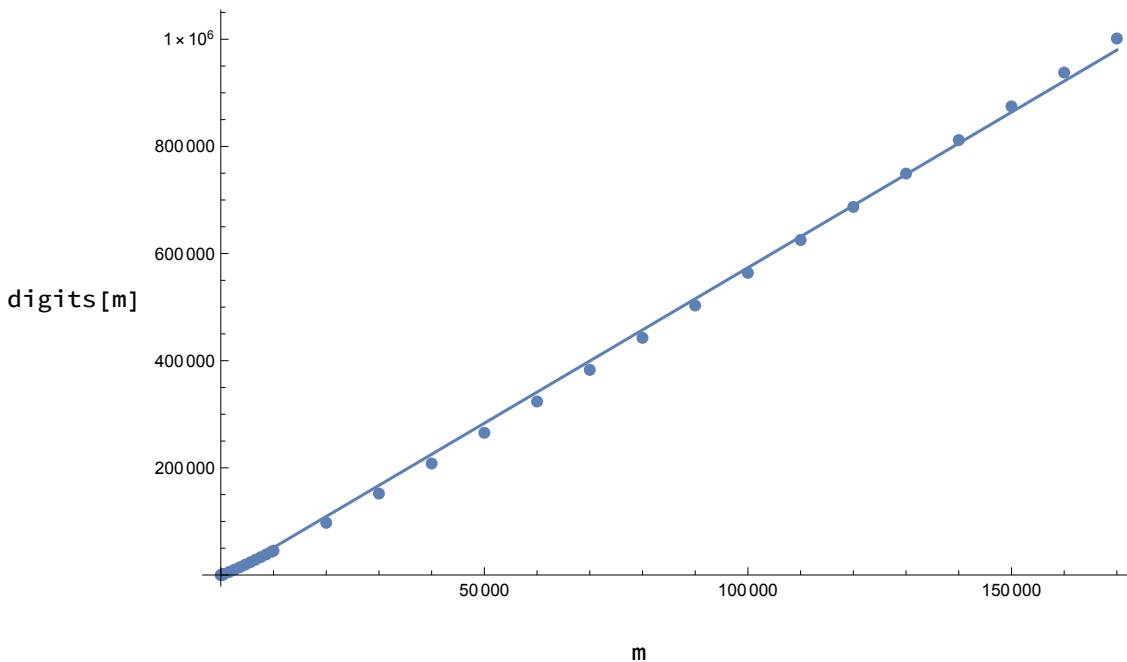
and

$$\text{PowerMod}[2, N[m, p] - 1, N[m, p]] = 1, \quad (10.)$$

where the latter is Fermat's little theorem. All of these methods require about the same computing time to determine that a number is not prime but Eq. 10 is much faster when the number is probably a prime (about the same time as for not a prime). However, for large primes the time to prove that a number is probably prime requires the use of a super computer although the total number of searches needed to find all of the primes in regions 1 and 2, the number of digits in large primorials, and the primorial itself can be easily computed with a desk top computer. In the remaining section the behavior of the numbers in regions 1 and 2 will be explored for small primes.

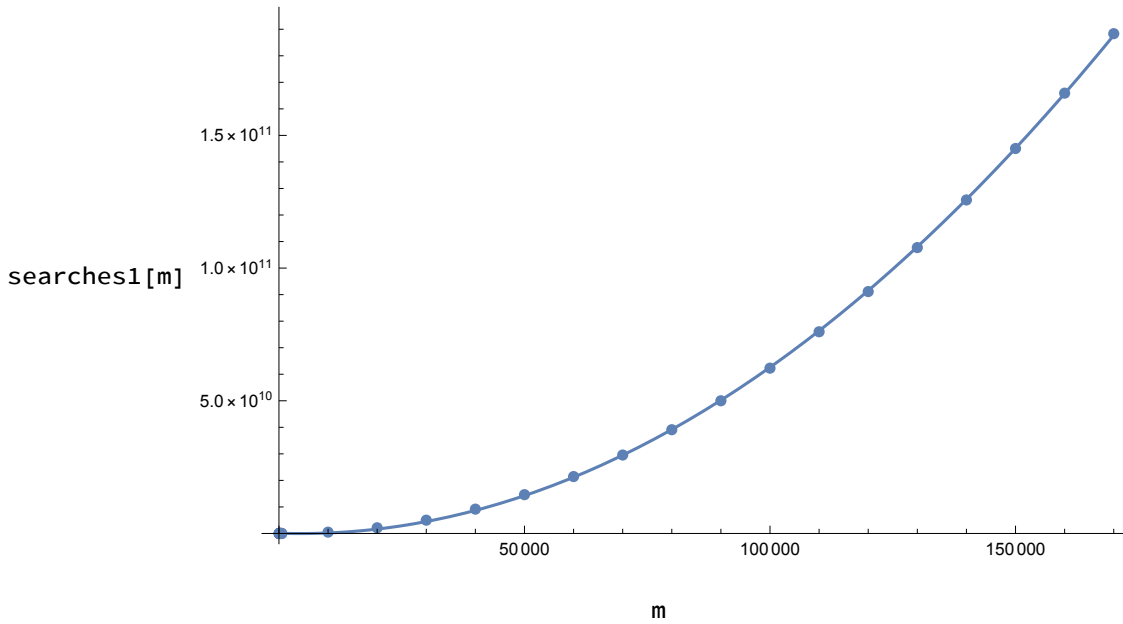
In Fig. 1 the dependence of the number of digits in the primorial as a function of m is shown which appears to be nearly a linear function of m .

Figure 1 : Number of digits in the m th primorial as a function of m (points). The solid line is the best fit straight line given as : $\text{digits}[m] = -6888.6 + 5.80435 m$.



In Fig. 2 the number of searches, $p(\text{max})[m]$, required to cover all of region 1 is shown as a function of m .

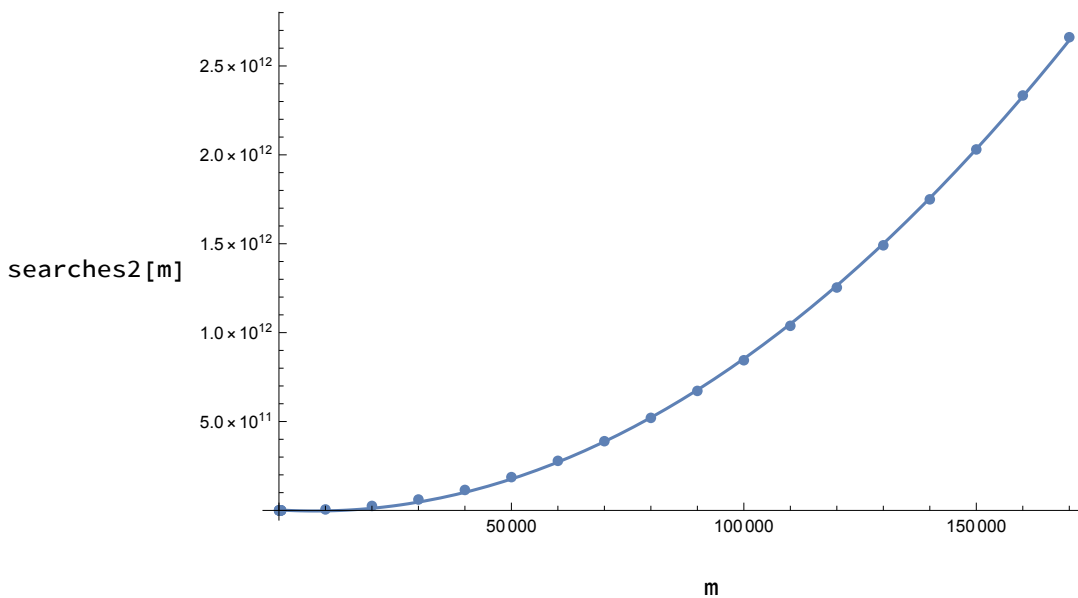
Figure 2 : Number of searches in region 1 required to find all possible primes of the type primorial plus or minus a prime as a function of m (points). The solid line is the best fit quadratic function given as : $\text{searches1}[m] = 9.52099 * 10^7 - 58716 m + 6.83784 m^2$



In Fig. 3 the same plot for the total number of searches required , $k(\max) [m]$, to find all primes in region 2 is shown with the best quadratic fit. Note that the number of searches required is more than an order of magnitude larger than for searches in region 1 for numbers with more than a million digits.

Figure 3 : Number of searches in region 2 required to find all possible primes of the type primorial plus an odd number as a function of m (points). The solid line is the best fit quadratic function given as :

$$\text{searches2}[m] = 2.74236 * 10^9 - 1.55020 * 10^6 m + 100.53557 m^2$$



The remaining quantities of interest : the number of primes in regions 1 and 2 as a function of m , the prime density as a function of m , and its reciprocal, the average number of searches per prime found , all involve determining probable primality which can only reasonably be determined using a desktop computer for small values of m . In Figures 4 - 7 the dependence of the total number of primes found in complete searches of regions 1 and 2 are shown along with best linear fits of the data for values of m up to 500.

Figure 4 : Total number of primes found in region 1 of the type primorial plus a prime number as a function of m (points). The solid line is the best fit to a linear function given as : $\text{prime1p}[m] = -162.3895 + 7.18096 m$

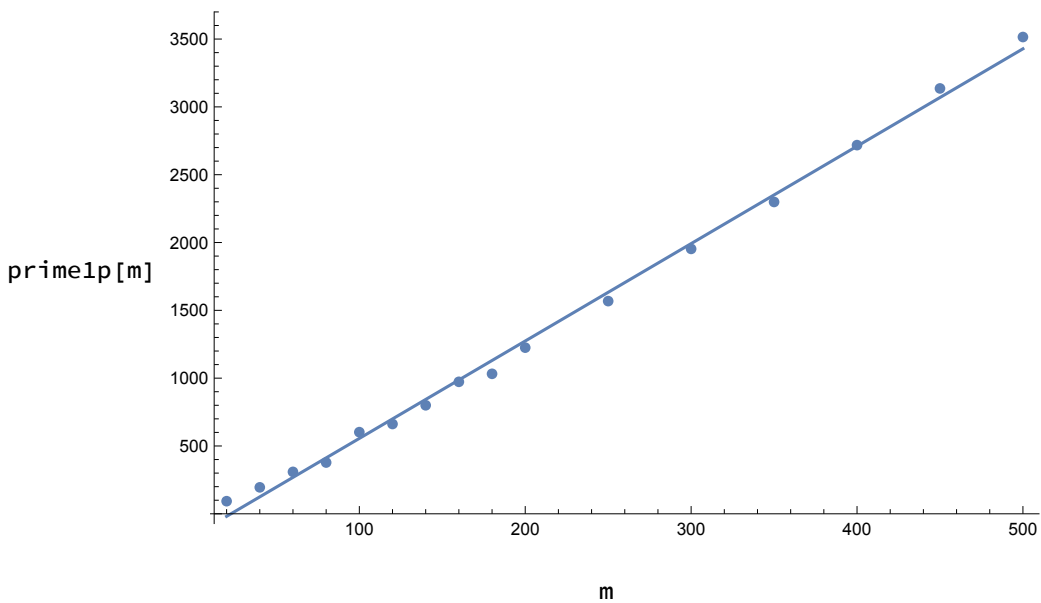


Figure 5 : Total number of primes found in region 1 of the type primorial minus a prime number as a function of m (points). The solid line is the best fit to a linear function given as : $\text{prime1m}[m] = -139.134797 + 6.99766 m$

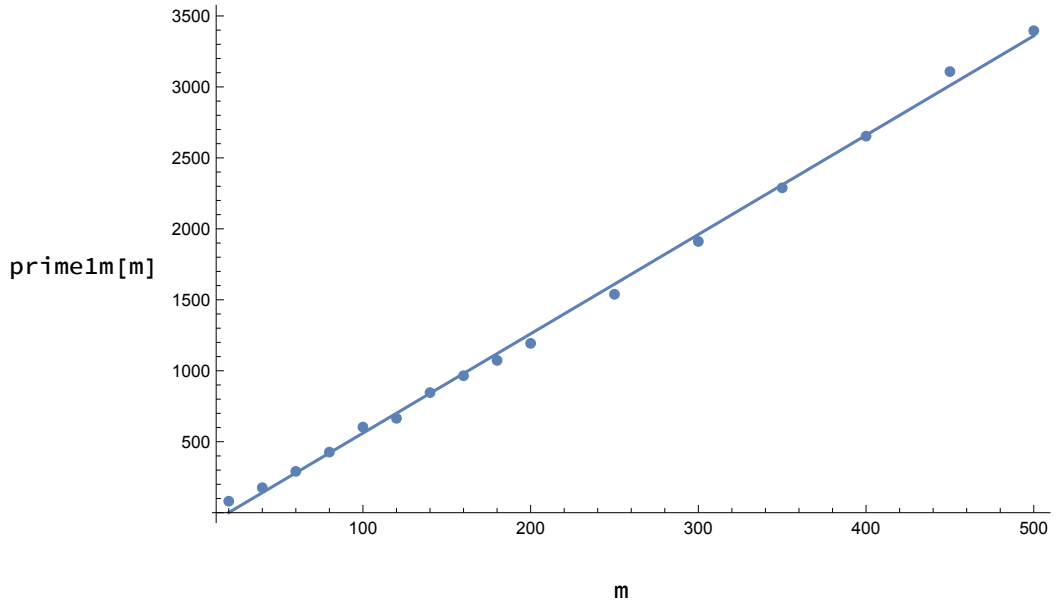


Figure 6 : Total number of primes found in region 2 of the type primorial plus all odd numbers as a function of m (points) . The solid line is the best fit to a linear function given as : $\text{prime2p}[m] = -178.9333 + 6.91721 m$

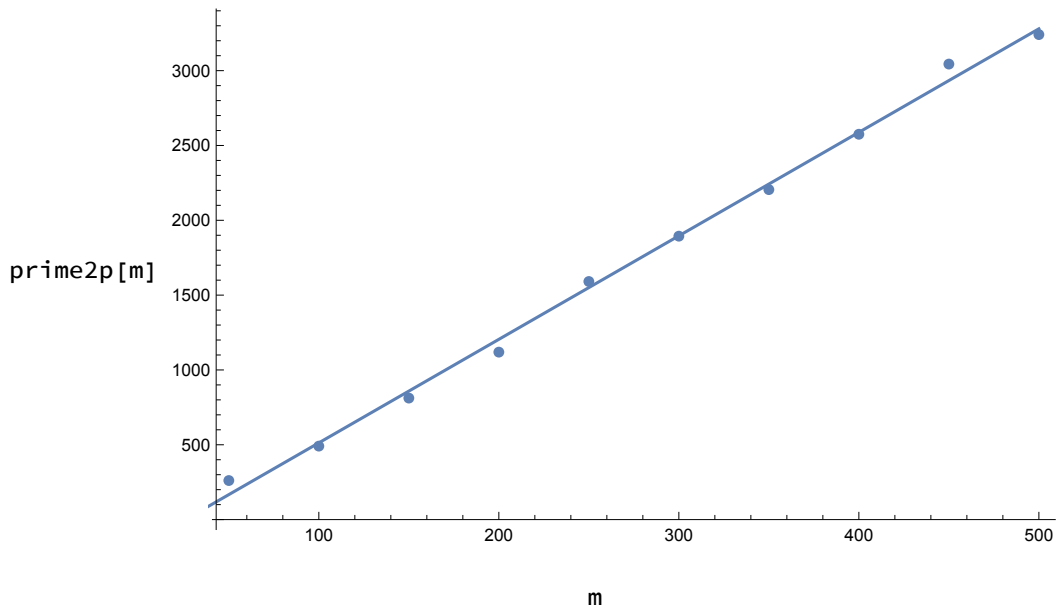
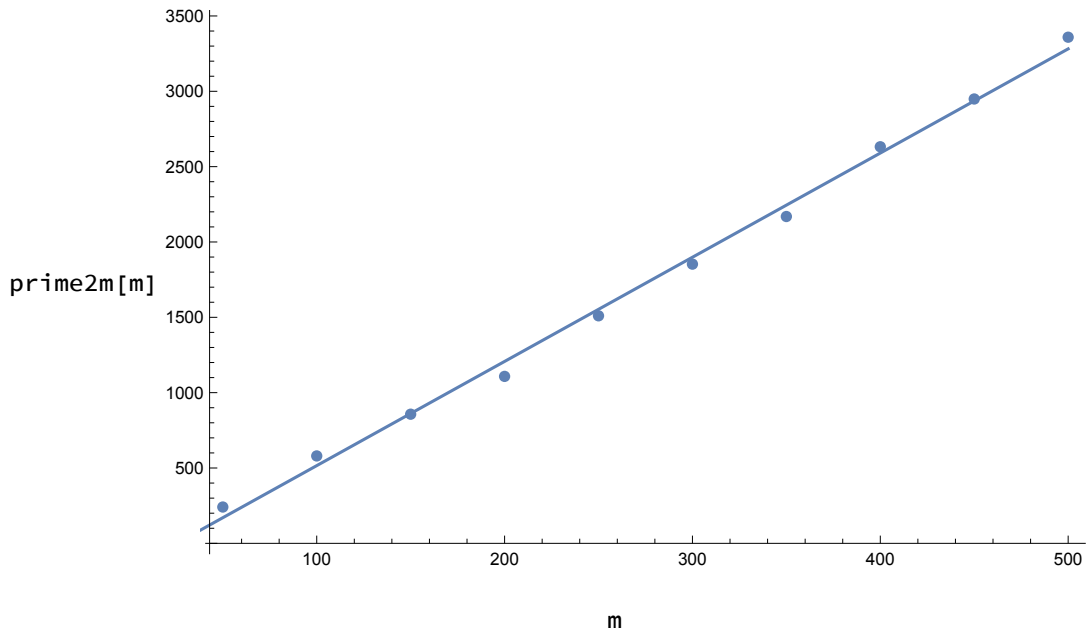


Figure 7 : Total number of primes found in region 2 of the type primorial minus all odd numbers as a function of m (points) . The solid line is the best fit to a linear function given as : $\text{prime2m}[m] = -175.73333 + 6.914667 m$



It should be noted that the total number of primes in region 1 is greater than the number in region 2 even if the numbers with $\text{Prime}[m]^2 + 2 * k$ not prime are included.

In Fig. 8 the average number of searches to find a primorial plus a prime in region 1 is displayed. Extrapolation of this result to $m = 170\,000$ (million digit primes) suggests that, on average, about 81 000 searches would be required to locate a probable prime of that size. Fig. 9 displays the results for a primorial minus a prime as a function of m with a projected number of searches to find a prime of about 84 000.

Figure 8 : Average number of searches to find a prime in region 1 of the type primorial plus a prime as a function of m (points). The solid line is the best fit to a linear function given as : $\text{search}/\text{primep}[m] = -1.67901 + 0.474315 m$

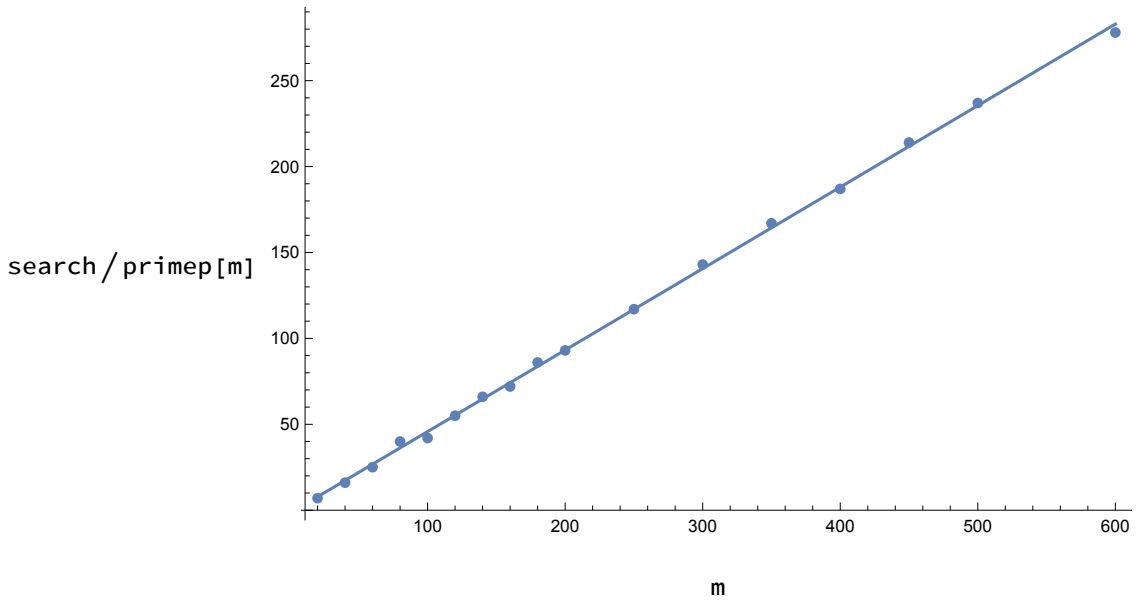
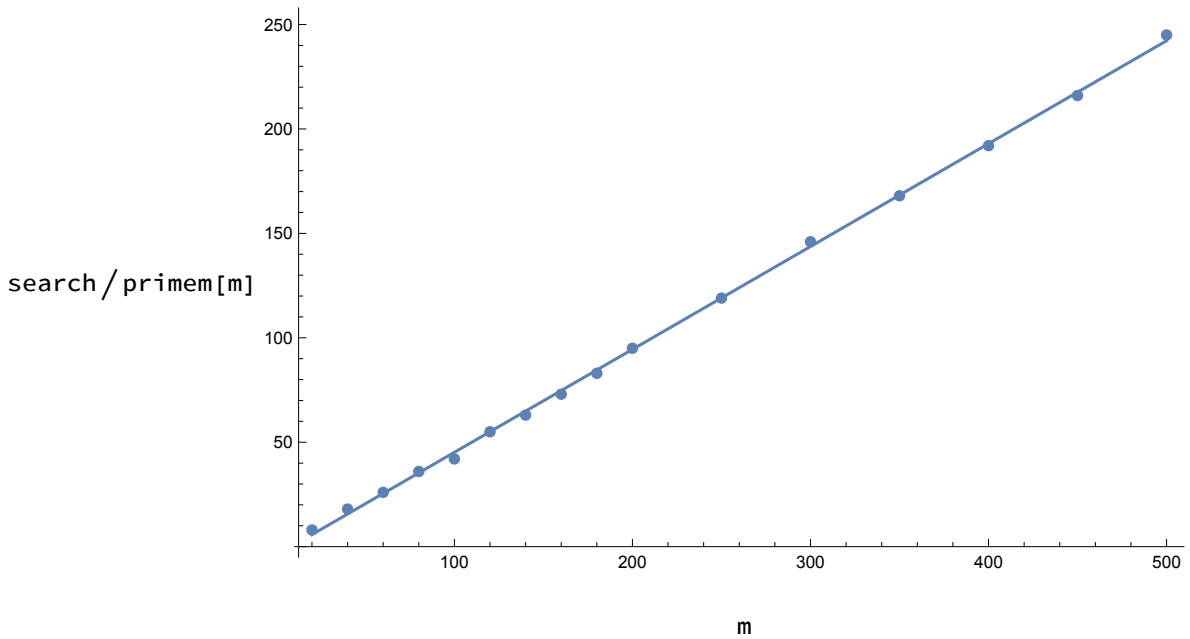


Figure 9 : Average number of searches to find a prime in region 1 of the type primorial minus a prime as a function of m (points). The solid line is the best fit to a linear function given as : $\text{search/primep}[m] = -4.08821 + 0.492660 m$



In Figures 10 and 11 the average number of searches to find a prime in region 2 are shown. For $m = 170\,000$ the extrapolated number of searches is more than $670\,000$ or approximately 8 times the number of searches required for region 1.

Figure 10 : Average number of searches to find a prime in region 2 of the type primorial plus all odd numbers as a function of m (points). The solid line is the best fit to a linear function given as : $\text{search}/\text{primep2}[m] = -109.19498 + 3.948692 m$

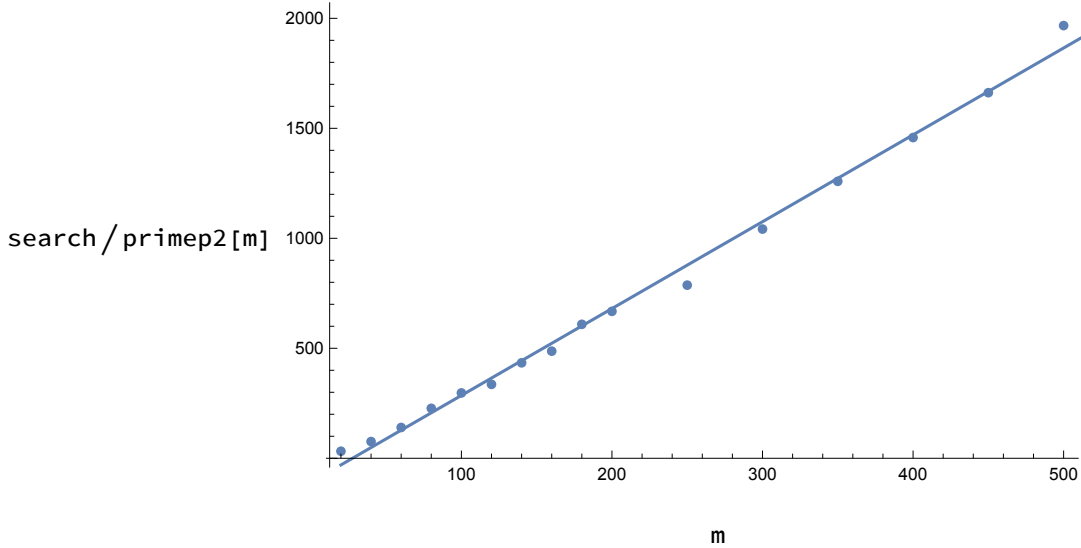
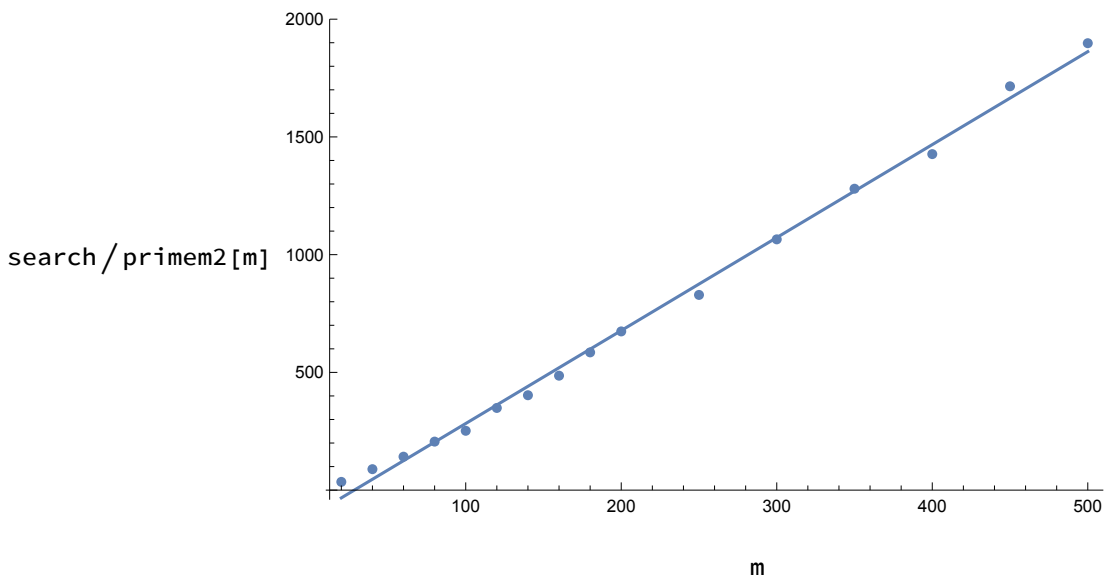


Figure 11 : Average number of searches to find a prime in region 2 of the type primorial minus all odd numbers as a function of m (points). The solid line is the best fit to a linear function given as : $\text{search}/\text{primem2}[m] = -111.94147 + 3.948079 m$



In summary, it has been shown that the number of searches required to find all primes in regions 1 and 2 is approximately a quadratic

function of m where m is the m th primorial. The number of digits in the primorial the number of probable primes in regions 1 and 2 , and the average number of searches required to find a probable prime in regions 1 and 2 are all approximately linear functions of m . 